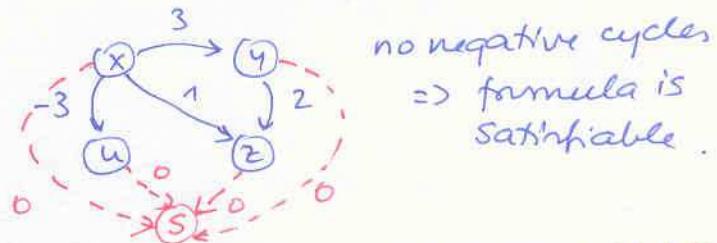


## DIFFERENCE LOGIC : Examples

Check the satisfiability of the following formulae in (positive) difference logic (in  $\mathbb{Q}$ ). In case of unsatisfiability find a satisfying assignment using the method presented in the lecture.

$$(1) \quad x-y \leq 3 \wedge y-z \leq 2 \wedge x-z \leq 1 \wedge x-u \leq -3$$

Associated graph :



no negative cycles  
=> formula is  
satisfiable.

To compute a satisfying assignment we add a new vertex  $s$  and edges with weight 0 from any node to  $s$ .

let  $\delta_{xs}$  be the length of the shortest path from vertex  $x$  to  $s$ .

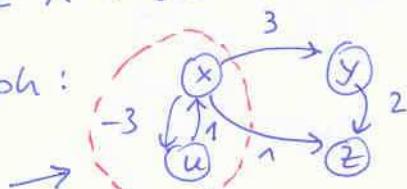
	$\delta_{ys}$	$\delta_{zs}$	$\delta_{us}$	$\delta_{ys}$	$\delta_{zs}$	$\delta_{us}$	$y$	$z$	$u$
$\delta_{ys}$	-	-	-	-	-	-	-	-	-
$\delta_{zs}$	-	-	-	-	-	-	-	-	-
$\delta_{us}$	-	-	-	-	-	-	-	-	-

Then  $\beta: X \rightarrow \mathbb{Q}$  defined by  $\beta(x) = \delta_{xs}$  for every  $x \in X$  is a satisfying assignment.

$\delta_{xs} = -3, \delta_{ys} = 0, \delta_{zs} = 0, \delta_{us} = 0 \Rightarrow \beta: X \rightarrow \mathbb{Q}$  defined by  $\beta(x) = -3, \beta(y) = \beta(z) = \beta(u) = 0$  is a satisfying assignment.

$$(2) \quad x-y \leq 3 \wedge y-z \leq 2 \wedge x-z \leq 1 \wedge x-u \leq -3 \wedge u-x \leq 1$$

Associated graph :



cycle of  
negative  
total weight

$\Rightarrow$  the formula is unsatisfiable.