# Universität Koblenz-Landau

### FB 4 Informatik

### Prof. Dr. Viorica Sofronie-Stokkermans

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# Exercises for "Decision Procedures for Verification" Exercise sheet 1

### Exercise 1.1: (5 P)

Determine which of the following formulas are valid/satisfiable/unsatisfiable:

- $(1) (P \wedge Q) \rightarrow (P \vee Q)$
- $(2) (P \lor Q) \to (P \land Q)$
- $(3) \neg (P \land \neg \neg P)$
- (4)  $Q \rightarrow \neg Q$
- (5)  $Q \wedge \neg Q$
- (6)  $\neg (\neg P \lor \neg \neg P)$
- $(7) ((P \to Q) \land (\neg P \to R)) \to (Q \lor R)$

#### Exercise 1.2: (5 P)

If F and G are propositional formulae then prove that the following are equivalent:

- (a)  $F \models G$ ;
- (b)  $\models F \rightarrow G$  (i.e.  $F \rightarrow G$  is valid);
- (c)  $F \wedge \neg G$  is unsatisfiable.

# Exercise 1.3: (2 P)

Prove: If N is a set of propositional formulas, then  $N \models F$  if and only if  $N \cup \{\neg F\}$  is unsatisfiable.

(A set of propositional formulas is unsatisfiable, if and only if for every valuation  $\mathcal{A}$  there is a formula G in the set such that  $\mathcal{A} \not\models G$ .)

## **Exercise 1.4:** (5 P)

Prove:

(1) If  $F_1, \ldots, F_n, G$  are propositional formulae then  $F_1 \wedge \cdots \wedge F_n \to G$  is valid iff every valuation which is a model of all the formulae  $F_1, F_2, \ldots F_n$  is also a model of G.

- (2) If  $F_1, \ldots, F_n, G$  are propositional formulae then the following are equivalent:
  - (a)  $\{F_1, ..., F_n\} \models G$
  - (b)  $F_1 \wedge \cdots \wedge F_n \to G$  is valid
  - (c)  $F_1 \wedge \cdots \wedge F_n \wedge \neg G$  is unsatisfiable.

## Supplementary exercises

### Exercise 1.5: (2 P)

Consider the formulae  $F_n = \bigvee_{i=1}^n (Q_i \wedge R_i)$  for  $n \in \mathbb{N}$ .

As a function of n, how many clauses are in:

- (1) the CNF formula F' constructed using the distributivity of disjunctions over conjunctions?
- (2) the CNF formula F'' obtained using the satisfiability-preserving translation to clause form?
- (3) For which n is the first approach better?

You will be able to solve 1.5(2) and (3) only after the satisfiability-preserving translation to clauses normal form is presented in the lecture on Monday, 31.10.2016.

Please submit your solution until Wednesday, November 2, 2016 at 13:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name on your solution.

Submission possibilities:

- By e-mail to sofronie@uni-koblenz.de with the keyword "Homework DP" in the subject.
- Put it in the box in front of Room B 222.