

Exercises for “Decision Procedures for Verification” Exercise sheet 11

Exercise 11.1: (4 P)

Use the Nelson-Oppen procedure for checking the satisfiability of the following formulae:

1. $1 \leq c \wedge c \leq 3 \wedge f(c) \not\approx f(1) \wedge f(c) \not\approx f(3) \wedge f(1) \not\approx f(2)$
in the combination $LI(\mathbb{Z}) \cup UIF_{\{f\}}$.
2. $f(c) \approx f(c+d) \wedge 1 \leq c \wedge c \leq d+e \wedge c+e \leq d \wedge d=1 \wedge f(c) \not\approx f(2)$
in the combination $LI(\mathbb{Z}) \cup UIF_{\{f\}}$.

Exercise 11.2: (4p P)

Let $\mathcal{T} = LI(\mathbb{Q})$, and let $Q := y \leq 1, R := x \leq y, P := y + y \leq 2, S := x \geq 1$. Use a DPLL(\mathcal{T}) method to check the satisfiability w.r.t. \mathcal{T} of the following set of clauses:

- (1) $\neg R \vee P$
- (2) $\neg Q \vee \neg P$
- (3) $R \vee P$
- (4) S

For checking the satisfiability of conjunctions of inequalities in $LI(\mathbb{Q})$ use the Fourier-Motzkin method.

In what follows we consider the theory of arrays defined in the lecture. We assume that the theory of indices \mathcal{T}_i is $LI(\mathbb{Z})$, and the theory of elements \mathcal{T}_e is $LI(\mathbb{Q})$.

Exercise 11.3: (2 P)

Which of the formulae below are (equivalent to formulae) in the array property fragment and which are not?

Justify your answer. (The universally quantified variables i, j are sort *index*; the indices k, l which are not universally quantified are considered to be constants of sort *index*)

- (1) $\forall i (a[i+1] > a[i])$
- (2) $\forall i (i < a[k] \rightarrow a[i] = a[k])$
- (3) $\forall i, j (l_1 \leq i \leq u_1 < l_2 \leq j \leq u_2 \rightarrow a[i] \leq a[j])$
- (3) $\forall i, j (l_1 < i \leq u_1 < l_2 \leq j \leq u_2 \rightarrow a[i] \leq a[j])$.

Supplementary exercises:

Exercise 11.4: (5 P)

Let \mathcal{T} be a theory with signature Σ and $\text{Mod}(\mathcal{T})$ be its class of models.

Show that if $\text{Mod}(\mathcal{T})$ is closed under products then \mathcal{T} is convex.

Exercise 11.5: (5 P)

We say that a theory \mathcal{T} is *stably infinite* if for every quantifier-free formula ϕ , ϕ is satisfiable in \mathcal{T} iff ϕ is satisfiable in a (countably) infinite model of \mathcal{T} .

Let $\mathcal{T}_1, \mathcal{T}_2$ be stably infinite theories with disjoint signatures. Prove that their combination $\mathcal{T}_1 \cup \mathcal{T}_2$ is stably infinite.

Please submit your solution until Wednesday, February 15, 2017 at 13:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name on your solution.

Submission possibilities:

- By e-mail to sofronie@uni-koblenz.de with the keyword “Homework DP” in the subject.
- Put it in the box in front of Room B 222.