

Let  $\mathcal{T}$  be the combination of  $LI(\mathbb{Q})$  (linear arithmetic over  $\mathbb{Q}$ ) and  $UIF_\Sigma$ , the theory of uninterpreted function symbols in a signature  $\Sigma$  containing the unary function  $h$ .

Check the satisfiability of the following ground formula w.r.t.  $\mathcal{T}$  using the deterministic version of the Nelson-Oppen procedure (after purifying the formulae check for entailment of equalities between shared constants and propagate the entailed equalities):

- $\phi : a + h(h(h(b)) + c) \approx e \wedge h(b) \approx b' \wedge e \approx h(b') + c \wedge a + h(e) \not\approx e$ .

### Solution

#### Step 1: Purification

$$a + h(\underbrace{h(h(b))}_{b_1} + c) \approx e \wedge h(b) \approx b' \wedge e \approx \underbrace{h(b')}_{b''} + c \wedge a + \underbrace{h(e)}_{e_1} \not\approx e.$$

$$\underbrace{\underbrace{b_1}_{c_1}}_{b_2}$$

$LI(\mathbb{Q})$	$UIF$
$a + b_2 \approx e$	$h(h(b)) \approx b_1$
$b_1 + c \approx c_1$	$h(c_1) \approx b_2$
$e \approx b'' + c$	$h(b) \approx b'$
$a + e_1 \not\approx e$	$h(b') \approx b''$
	$h(e) \approx e_1$
$\phi_1$	$\phi_2$

Both  $\phi_1$  and  $\phi_2$  are satisfiable, so we start Step 2.

#### Step 2: Propagation

Shared constants:  $\{b_1, b_2, b'', e, e_1\}$ .

We notice that  $\phi_2 \models_{UIF} b_1 \approx b''$

Indeed, as  $b(b) \approx b'$  we have  $h(h(b)) \approx h(b')$ .  $\phi_1$  contains  $h(h(b)) \approx b_1 \wedge h(b') \approx b''$ , thus  $\phi_2 \models_{UIF} b_1 \approx b''$ .

We propagate  $b_1 \approx b''$  and obtain:

$LI(\mathbb{Q})$	$UIF$
$a + b_2 \approx e$	$h(h(b)) \approx b_1$
$b_1 + c \approx c_1$	$h(c_1) \approx b_2$
$e \approx b'' + c$	$h(b) \approx b'$
$a + e_1 \not\approx e$	$h(b') \approx b''$
	$h(e) \approx e_1$
$b_1 \approx b''$	$b_1 \approx b''$
$\phi_1^1$	$\phi_2^1$

We notice that  $\phi_1^1 \models_{LI(Q)} e \approx c_1$

Inded, as  $b_1 \approx b''$  we have  $b_1 + c \approx b'' + c$ , so since  $\phi_1^1$  contains  $b_1 + c \approx c_1$  and  $b'' + c \approx e$  we obtain  $\phi_1^1 \models e \approx c_1$ .

We propagate  $e \approx c_1$  and obtain:

$LI(Q)$	$UIF$
$a + b_2 \approx e$	$h(h(b)) \approx b_1$
$b_1 + c \approx c_1$	$h(c_1) \approx b_2$
$e \approx b'' + c$	$h(b) \approx b'$
$a + e_1 \not\approx e$	$h(b') \approx b''$
	$h(e) \approx e_1$
$b_1 \approx b''$	$b_1 \approx b''$
$e \approx c_1$	$e \approx c_1$
$\phi_1^2$	$\phi_2^2$

We notice that  $\phi_2^2 \models_{UIF} b_2 \approx e_1$

Inded, as  $e \approx c_1$  we have  $h(e) \approx h(c_1)$ .  $\phi_2^2$  contains  $h(e) \approx e_1 \wedge h(c_1) \approx b_2$ , thus  $\phi_2 \models_{UIF} b_2 \approx e_1$ .

We propagate  $b_2 \approx e_1$  and obtain:

$LI(Q)$	$UIF$
$a + b_2 \approx e$	$h(h(b)) \approx b_1$
$b_1 + c \approx c_1$	$h(c_1) \approx b_2$
$e \approx b'' + c$	$h(b) \approx b'$
$a + e_1 \not\approx e$	$h(b') \approx b''$
	$h(e) \approx e_1$
$b_1 \approx b''$	$b_1 \approx b''$
$e \approx c_1$	$e \approx c_1$
$b_2 \approx e_1$	$b_2 \approx e_1$
$\phi_1^3$	$\phi_2^3$

We now show that  $\phi_1^3$  is unsatisfiable:

From  $b_2 \approx e_1$  it follows that  $a + b_2 \approx a + e_1$ .  
But  $\phi_1^3$  contains  $a + b_2 \approx e$  and  $a + e_1 \not\approx e$ . Contradiction.

Thus, the formula  $\phi$  is unsatisfiable.