

Exercises for “Decision Procedures for Verification” Exercise sheet 8

Exercise 8.1: (2 P)

Let ϕ be the following (ground) formula:

$$f(f(c)) \approx f(c) \wedge f(f(c)) \approx f(d) \wedge d \not\approx f(c).$$

- (1) Compute $FLAT(\phi)$ (the formula obtained by recursively replacing, in a bottom-up fashion, any term of the form $f(c')$, where c' is a constant, with a new constant).
- (2) Compute $FC(\phi)$ (the set of functional consistency axioms associated with the flattening above):

$$FC(\phi) = \{c_1 \approx c_2 \rightarrow d_1 \approx d_2 \mid d_i \text{ is introduced as an abbreviation for } f(c_i)\}.$$

- (3) Check whether $FLAT(\phi) \wedge FC(\phi)$ is satisfiable.
- (4) Is ϕ is satisfiable? Justify your answer.

Exercise 8.2: (6 P)

Check the satisfiability of the following ground formulae using the algorithm based on congruence closure presented in the lecture on Thursday, 20.12.2018 (set of slides presented on 17.12 and 20.12.2018).

- (1) $\phi_1 = f(f(c)) \approx f(c) \wedge f(f(c)) \approx f(d) \wedge d \not\approx f(c)$.
- (2) $\phi_2 = f(f(c)) \approx f(c) \wedge f(c) \approx d \wedge f(d) \not\approx f(f(c))$.

Supplementary exercise:

Exercise 8.3: (3 P)

Prove the \Rightarrow part in the correctness proof of the algorithm for checking the validity of a conjunction of literals in UIF, under the assumption that an algorithm for computing the congruence closure of a set R of pairs of vertices in a graph G exists.

Let $\phi := \bigwedge_{i=1}^n s_i \approx t_i \wedge \bigwedge_{j=1}^m s'_j \not\approx t'_j$ be a ground formula. Let $G = (V, E)$ be the labelled directed graph constructed from ϕ as in the description of the congruence closure algorithm based on Union/Find. Let $R = \{(v_{s_i}, v_{t_i}) \mid i \in \{1, \dots, n\}\}$, and let R^c be the congruence closure of R .

- (1) \mathcal{A} is a Σ -structure such that $\mathcal{A} \models \phi$. Prove that $[v_s]_{R^c} = [v_t]_{R^c}$ implies that $\mathcal{A} \models s = t$.

(2) Assume that ϕ is satisfiable. Prove that $[v_{s'_j}]_{R^c} \neq [v_{t'_j}]_{R^c}$.

Hint: Use the fact that if $[v_s]_{R^c} = [v_t]_{R^c}$ then there is a derivation for $(v_s, v_t) \in R^c$ in the calculus defined before; use induction on the length of derivation to show that $\mathcal{A} \models s = t$.

Please submit your solution until Wednesday, January 9, 2018 at 12:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name on your solution.

Submission possibilities:

- By e-mail to sofronie@uni-koblenz.de with the keyword “Homework DP” in the subject.
- Put it in the box in front of Room B 222.