

Exercise 5.1

$\Sigma = (\Omega, \Pi)$ with $\Omega = \{b, f\}$ and $\Pi = \{P\}$.

(1) How many Herbrand interpretations over Σ exist?

Solution: $\mathcal{H} = (U_{\mathcal{H}}, \{f_{\mathcal{H}}, b_{\mathcal{H}}\}, \{P_{\mathcal{H}}\})$ is a Herbrand interpretation

over Σ iff (1) $U_{\mathcal{H}} = T_{\Sigma} = \{b, f(b), f^2(b), \dots\}$

(2) $f_{\mathcal{H}} : T_{\Sigma} \rightarrow T_{\Sigma}$ is defined by $f_{\mathcal{H}}(f^n(b)) = f(f^n(b)) = f^{n+1}(b)$.

(3) $b_{\mathcal{H}} = b \in T_{\Sigma}$

(3) $P_{\mathcal{H}} \subseteq T_{\Sigma}$

The universe and the interpretation of the function symbols are fixed. Therefore, there are as many Herbrand interpretations over Σ as subsets of T_{Σ} .

$$|\{\mathcal{H} \mid \mathcal{H} \text{ Herbrand interpretation over } \Sigma\}| = |\mathcal{P}(T_{\Sigma})| = |\mathcal{P}(\mathbb{N})|$$

Since $\mathcal{P}(\mathbb{N})$ is not countable, the set of all Herbrand interpretations over Σ is not countable.

because there is a bijection from \mathbb{N} to T_{Σ}
 $i: \mathbb{N} \rightarrow T_{\Sigma}, i(n) = f^n(b)$.

[It is known, in fact, that $|\mathcal{P}(\mathbb{N})| = |\mathbb{R}|$. The proof is given in a different file]

(2) Consider $F := P(f(f(b))) \wedge \forall x (P(x) \rightarrow P(f(x)))$.
 How many different Herbrand models does F have?

Let $\mathcal{A} = (T_{\Sigma}, \{f_{\mathcal{A}} : T_{\Sigma} \rightarrow T_{\Sigma}, f_{\mathcal{A}}(f^n(b)) = f^{n+1}(b), b_{\mathcal{A}} = b \in T_{\Sigma}\}, P_{\mathcal{A}} \subseteq T_{\Sigma})$ be a Herbrand interpretation such that $\mathcal{A} \models F$.

Then $\mathcal{A} \models P(f(f(b)))$, i.e. $P_{\mathcal{A}}(f(f(b))) = 1$ (i.e. $f(f(b)) \in P_{\mathcal{A}}$) (*)

$\mathcal{A} \models \forall x (P(x) \rightarrow P(f(x)))$, i.e. $\forall \beta: X \rightarrow T_{\Sigma}, \mathcal{A}(\beta) (\forall x P(x) \rightarrow P(f(x))) = 1$
 i.e. $\forall \beta: X \rightarrow T_{\Sigma} \min_{t \in T_{\Sigma}} (\mathcal{A}(\beta[x \mapsto t]) (P(x) \rightarrow P(f(x)))) = 1$

i.e. $\forall \beta: X \rightarrow T_{\Sigma}$ and $\forall t \in T_{\Sigma} : \mathcal{A}(\beta[x \mapsto t]) (P(x) \rightarrow P(f(x))) = 1$
 i.e. $\forall \beta: X \rightarrow T_{\Sigma}$ and $\forall t \in T_{\Sigma} : P_{\mathcal{A}}(t) \rightarrow P_{\mathcal{A}}(f(t)) = 1$
 i.e. $\forall \beta: X \rightarrow T_{\Sigma}$ and $\forall t = f^n(b) : P_{\mathcal{A}}(f^n(b)) \rightarrow P_{\mathcal{A}}(f^{n+1}(b)) = 1$ (**)

Hence, if $\mathcal{A} \models F$ then

(i) $P_{\mathcal{A}}(f(f(b))) = 1$ (*)

(ii) $(P_{\mathcal{A}}(f(f(b))) \rightarrow P_{\mathcal{A}}(f(f(f(b)))) = 1$ (**)

(iii) $(P_{\mathcal{A}}(f^3(b)) \rightarrow P_{\mathcal{A}}(f^4(b))) = 1$ (**)

hence (iii) $P_{\mathcal{A}}(f^3(b)) = 1$

hence (iv) $P_{\mathcal{A}}(f^4(b)) = 1$

...
 one can prove by induction that $P_{\mathcal{A}}(f^n(b)) = 1$ for all $n \geq 2$. Nothing is known about $P_{\mathcal{A}}(b), P_{\mathcal{A}}(f(b))$

We showed that if A is a Heibrand interpretation such that $A \models F$, then $P_A(f^n(b)) = 1$ for all $n \geq 2$.

Case 1 $P_A(b) = 1$. Then by $(*)$ $P_A(f(b)) = 1$, hence $P_A = \{b, f(b), f^2(b), \dots\} = \bar{\Sigma}$.

Case 2 $P_A(b) = 0$. We distinguish two

Subcases: Case 2.1: $P_A(b) = 1$. Then $P_A = \{f(b), f^2(b), \dots\} = \bar{\Sigma} \setminus \{b\}$

Case 2.2 $P_A(b) = 0$. Then $P_A = \{f^2(b), f^3(b), \dots\} = \bar{\Sigma} \setminus \{b, f(b)\}$.

There are 3 interpretations:

$$U_1 = (\bar{\Sigma}, \{f_{U_1}, b_{U_1}\}, P_{U_1} = \bar{\Sigma})$$

$$U_2 = (\bar{\Sigma}, \{f_{U_2}, b_{U_2}\}, P_{U_2} = \{f^n(b) \mid n \geq 1\})$$

$$U_3 = (\bar{\Sigma}, \{f_{U_3}, b_{U_3}\}, P_{U_3} = \{f^n(b) \mid n \geq 2\}).$$

Standard,
in any Heibrand int.

(3) Every Heibrand interpretation which is a model of F is also a model of $G = \forall x p(f(f(x)))$.

Give an example of an algebra that is a model of F but not of G .

Solution: let $A = (U_A, \{f_A, b_A\}, P_A)$ where

$$U_A = \{a, b, f(a), f(b), f^2(a), f^2(b), \dots\}$$

$$= \{f^n(a) \mid n \geq 0\} \cup \{f^n(b) \mid n \geq 0\}$$

$$f_A: U_A \rightarrow U_A \text{ defined by } f_A(f^n(a)) = f^{n+1}(a)$$

$$f_A(f^n(b)) = f^{n+1}(b).$$

$$b_A \in U_A \quad b_A = b.$$

$$P_A = \{f^n(b) \mid n \geq 2\}.$$

Then $A \models F$ but $A \not\models G$: *the counter example is...*

In fact, for all $\beta: X \rightarrow U_A$: $A(\beta)(G) = 0$:

$$A(\beta)(G) = A(\beta)(\forall x p(f(f(x)))) = \min_{t \in U_A} P_A(f(f(t))) = 0$$

because for $t = \underline{a}$ we have $P_A(f^2(a)) = 0$

(4) Let A be a Herbrand interpretation over Σ and \sim the binary relation on T_Σ defined by:

$$t_1 \sim t_2 \text{ iff } (\forall x f^3(x) \approx x) \models t_1 \approx t_2$$

a) Is \sim a congruence relation on A ?

Solution

\sim is a congruence relation on A iff $\left[\begin{array}{l} \sim \subseteq T_\Sigma \times T_\Sigma \text{ is an equivalence relation} \\ \text{and it is "compatible" with } \neq, \neq. \end{array} \right.$

\sim is an equivalence relation

(i) \sim is reflexive, i.e. $\forall t \in T_\Sigma, t \sim t$.

Proof: We show that $\forall x f^3(x) \approx x \models t \approx t$, i.e. that

for all A' Σ -structure, $\beta: X \rightarrow U_{A'}$, if $A'(\beta)(\forall x f^3(x) \approx x) = 1$ then

then $A'(\beta)(t) = A'(\beta)(t)$,

always true, q.e.d.

always equal

(ii) \sim is symmetric, i.e. $\forall t \in T_\Sigma$, if $t_1 \sim t_2$ then $t_2 \sim t_1$.

Proof: Assume $t_1 \sim t_2$. We show that $t_2 \sim t_1$, i.e. $(\forall x f^3(x) \approx x) \models t_2 \approx t_1$.

Let A' be a Σ -structure and $\beta: X \rightarrow U_{A'}$ such that $A'(\beta)(\forall x f^3(x) \approx x) = 1$.

Since we assumed $t_1 \sim t_2$, i.e. $(\forall x f^3(x) \approx x) \models t_1 \approx t_2$, we know that

$A'(\beta)(t_1) = A'(\beta)(t_2)$. But then $A'(\beta)(t_2) = A'(\beta)(t_1)$.

(iii) \sim is transitive, i.e. $\forall t \in T_\Sigma$ if $t_1 \sim t_2$ and $t_2 \sim t_3$ then $t_1 \sim t_3$.

Proof: Assume $t_1 \sim t_2$ and $t_2 \sim t_3$.

We prove that $t_1 \sim t_3$ i.e. $(\forall x f^3(x) \approx x) \models t_1 \approx t_3$.

Let A' be a Σ -structure and $\beta: X \rightarrow U_{A'}$ be such that $A'(\beta)(\forall x f^3(x) \approx x) = 1$.

- Since we assumed that $t_1 \sim t_2$, i.e. $(\forall x f^3(x) \approx x) \models t_1 \approx t_2$, we know that $A'(\beta)(t_1) = A'(\beta)(t_2)$.

- Since we assumed that $t_2 \sim t_3$ i.e. $(\forall x f^3(x) \approx x) \models t_2 \approx t_3$, we know that $A'(\beta)(t_2) = A'(\beta)(t_3)$.

It follows that $A'(\beta)(t_1) = A'(\beta)(t_2) = A'(\beta)(t_3)$, hence

$$A'(\beta)(t_1) = A'(\beta)(t_3).$$

thus We showed that $(\forall x f^3(x) \approx x) \models t_1 \approx t_3$, i.e. that $t_1 \sim t_3$.

(iv) \sim is "compatible" with f , i.e. $\forall t_1, t_2$ if $t_1 \sim t_2$ then $f(t_1) \sim f(t_2)$

Proof: Assume $t_1 \sim t_2$, i.e. $(\forall x f^3(x) = x) \models t_1 \approx t_2$ (*)

We show that $f(t_1) \sim f(t_2)$, i.e. $(\forall x f^3(x) = x) \models f(t_1) \approx f(t_2)$,

let \mathcal{A}' be a Σ -structure, $\beta: X \rightarrow U_{\mathcal{A}'}$ be st. $\mathcal{A}' \models \beta \models \forall x f^3(x) = x$.

By (*) we know that $\mathcal{A}' \models \beta \models (t_1) = (t_2)$.

Then $\mathcal{A}' \models \beta \models (f(t_1)) = (f(t_2))$ because f is a function symbol.

This shows that $(\forall x f^3(x) = x) \models f(t_1) \approx f(t_2)$, hence that $f(t_1) \sim f(t_2)$.

(v) we study compatibility with p .

We check whether $\forall t_1, t_2 \in T_{\Sigma}$ if $t_1 \sim t_2$ and $p_{\mathcal{A}}(t_1) = 1$ then $p_{\mathcal{A}}(t_2) = 1$

This is not always the case:

We can take for instance the Herbrand interpretation

$$\mathcal{A} = (T_{\Sigma}, \{f_{\mathcal{A}}, b_{\mathcal{A}}\}, P_{\mathcal{A}} = \{f^3(b)\})$$

usual definition

i.e. only $P_{\mathcal{A}}(f^3(b))$ is true,

For $t_1 = b$ and $t_2 = f^3(b)$ we have:

$t_1 \sim t_2$ and $p_{\mathcal{A}}(t_1) = 1$, but $p_{\mathcal{A}}(t_2) = 0$.

\sim is not a congruence on all Herbrand interpretations.

It is a congruence only for these Herbrand interpretations:

$$\mathcal{A} = (T_{\Sigma}, \{f_{\mathcal{A}}, b_{\mathcal{A}}\}, P_{\mathcal{A}}) \text{ where } P_{\mathcal{A}}(b) = P_{\mathcal{A}}(f^{3k}(b)) \text{ for all } k$$

$$\text{and } P_{\mathcal{A}}(f(b)) = P_{\mathcal{A}}(f^{3k+1}(b)) \text{ for all } k$$

$$\text{and } P_{\mathcal{A}}(f^2(b)) = P_{\mathcal{A}}(f^{3k+2}(b)) \text{ for all } k.$$

Assume that $\mathcal{A} = (T_{\Sigma}, \{f_A, g_A\}, P_A)$ has the property that

$$P_A(b) = P_A(f^{3k}(b))$$

$$P_A(f(b)) = P_A(f^{3k+1}(b))$$

$$P_A(f^2(b)) = P_A(f^{3k+2}(b)) \quad \text{for all } k \in \mathbb{N}.$$

We can build the quotient

$\mathcal{A}/\sim = (T_{\Sigma}/\sim, \{f_{\mathcal{A}/\sim}, g_{\mathcal{A}/\sim}\}, P_{\mathcal{A}/\sim})$ as follows:

$$T_{\Sigma}/\sim = \{ [b], [f(b)], [f^2(b)] \}$$

$$\text{where } [b] = \{ t \in T_{\Sigma} \mid t \sim b \} = \{ b, f^3(b), \dots, f^{3k}(b), \dots \}$$

$$[f(b)] = \{ t \in T_{\Sigma} \mid t \sim f(b) \} = \{ f(b), f^4(b), \dots, f^{3k+1}(b), \dots \}$$

$$[f^2(b)] = \{ t \in T_{\Sigma} \mid t \sim f^2(b) \} = \{ f^2(b), f^5(b), \dots, f^{3k+2}(b), \dots \}$$

are the equivalence classes wrt. \sim .

$$f_{\mathcal{A}/\sim}([b]) = [f(b)]$$

$$f_{\mathcal{A}/\sim}([f(b)]) = [f^2(b)]$$

$$f_{\mathcal{A}/\sim}([f^2(b)]) = [b].$$

$$P_{\mathcal{A}/\sim}([b]) := P_A(b)$$

$$P_{\mathcal{A}/\sim}([f(b)]) := P_A(f(b))$$

$$P_{\mathcal{A}/\sim}([f^2(b)]) := P_A(f(f(b))).$$

(well-defined because congruence)

(5)