

Non-classical logics

Lecture 10: Infinitely-valued logics (Part 1)

Viorica Sofronie-Stokkermans

sofronie@uni-koblenz.de

Łukasiewicz logics

Łukasiewicz logics

$$\mathcal{L}_n, n \in \mathbb{N} \quad W_n = \{0, \frac{1}{n-1}, \frac{2}{n-1}, \dots, 1\}$$

$$\mathcal{L}_{\mathbb{N}_0} \quad W_{\mathbb{N}_0} = [0, 1] \cap \mathbb{Q}$$

$$\mathcal{L}_{\mathbb{N}_1} \quad W_{\mathbb{N}_1} = [0, 1]$$

Logical operations: $\vee, \wedge, \neg, \Rightarrow$

- $\vee = \max$
- $\wedge = \min$
- $\neg x = 1 - x$
- $x \Rightarrow y = \min(1, 1 - x + y)$

Łukasiewicz logics

Łukasiewicz implication $x \Rightarrow_{\mathcal{L}_n} y = \min(1, 1 - x + y)$

\mathcal{L}_n

\Rightarrow	0	$\frac{1}{n-1}$	$\frac{2}{n-1}$	\dots	$\frac{n-2}{n-1}$	1
0	1	1	1	\dots	1	1
$\frac{1}{n-1}$	$\frac{n-2}{n-1}$	1	1	\dots	1	1
$\frac{2}{n-1}$	$\frac{n-3}{n-1}$	$\frac{n-2}{n-1}$	1	\dots	1	1
\dots						
1	0	$\frac{1}{n-1}$	$\frac{2}{n-1}$	\dots	$\frac{n-2}{n-1}$	1

Łukasiewicz logics

Theorems.

1. For $n, m \in \mathbb{N}$, s.t. $(m - 1) | (n - 1)$, we have
 $\text{Tautologies}(\mathcal{L}_n) \subseteq \text{Tautologies}(\mathcal{L}_m)$
2. $\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \text{Tautologies}(\mathcal{L}_{\aleph_1})$
3. $\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \bigcap \{ \text{Tautologies}(\mathcal{L}_n) \mid n \geq 2, n \in \mathbb{N} \}$

Proofs

Theorem. For $n, m \in \mathbb{N}$, s.t. $(m - 1) | (n - 1)$, we have
 $\text{Tautologies}(\mathcal{L}_n) \subseteq \text{Tautologies}(\mathcal{L}_m)$

Proof

Assume $(m - 1) | (n - 1)$. Then $W_m \subseteq W_n$. Assume $F \in \text{Tautologies}(\mathcal{L}_n)$. Then F evaluates to 1 under every valuation into W_n , hence also under every valuation into W_m , so $F \in \text{Tautologies}(\mathcal{L}_m)$

Proofs

Theorem. For $n, m \in \mathbb{N}$, s.t. $(m - 1) | (n - 1)$, we have
 $\text{Tautologies}(\mathcal{L}_n) \subseteq \text{Tautologies}(\mathcal{L}_m)$

Remark: the converse also holds

If $\text{Tautologies}(\mathcal{L}_n) \subseteq \text{Tautologies}(\mathcal{L}_m)$ then $(m - 1) | (n - 1)$.

(This will be discussed in the next exercise session.)

Proofs

Theorem.

$$\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \text{Tautologies}(\mathcal{L}_{\aleph_1})$$

Proof.

" \supseteq " : Since $[0, 1] \cap \mathbb{Q} \subseteq [0, 1]$, it is clear that

$$\text{Tautologies}(\mathcal{L}_{\aleph_1}) \subseteq \text{Tautologies}(\mathcal{L}_{\aleph_0})$$

Proofs

Theorem.

$$\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \text{Tautologies}(\mathcal{L}_{\aleph_1})$$

Proof.

" \subseteq " : Let $F \in \text{Tautologies}(\mathcal{L}_{\aleph_0})$. Then for every assignment of values in $[0, 1] \cap \mathbb{Q}$ to the propositional variables $\{P_1, \dots, P_n\}$ of F evaluates to 1.

We can associate a function $f_F : [0, 1]^n \rightarrow [0, 1]$ with F which is defined as follows:

For all $(x_1, \dots, x_n) \in [0, 1]^n$ let $\mathcal{A} : \{P_1, \dots, P_n\} \rightarrow [0, 1]$ be defined by $\mathcal{A}(P_i) = x_i$.

We define $f_F(x_1, \dots, x_n) := \mathcal{A}(F)$

It can be proved by structural induction that f_F is a continuous function.

Let $(a_1, \dots, a_n) \in [0, 1]^n$. It is now sufficient to choose sequences of rational numbers converging to a_1, \dots, a_n respectively. $f_F(a_1, \dots, a_n)$ is the limit of the sequence defined this way, hence its value is 1.

Proofs

$$\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \bigcap \{ \text{Tautologies}(\mathcal{L}_n) \mid n \geq 2, n \in \mathbb{N} \}$$

Proof.

" \subseteq " : Follows from the fact that $W_n \subseteq [0, 1] \cap \mathbb{Q}$ for every $n \in \mathbb{N}$.

Proofs

$$\text{Tautologies}(\mathcal{L}_{\aleph_0}) = \bigcap \{ \text{Tautologies}(\mathcal{L}_n) \mid n \geq 2, n \in \mathbb{N} \}$$

Proof. " \supseteq "

Let F be a formula with prop. variables $\{P_1, \dots, P_k\}$ s.t. $F \notin \text{Tautologies}(\mathcal{L}_{\aleph_0})$.
Then there exists $\mathcal{A} : \{P_1, \dots, P_k\} \rightarrow [0, 1] \cap \mathbb{Q}$ s.t. $\mathcal{A}(F) \neq 1$.

Assume that $\mathcal{A}(P_1) = \frac{q_1}{p_1}, \dots, \mathcal{A}(P_k) = \frac{q_k}{p_k}$

Let $m = \text{lcm}(p_1, \dots, p_k)$. Then it is easy to see that $\mathcal{A}(P_i) \in W_{m+1}$ for all $1 \leq i \leq k$.

We thus constructed a valuation $\mathcal{A} : \{P_1, \dots, P_k\} \rightarrow W_m$ such that $\mathcal{A}(F) \neq 1$.
Hence, $F \notin \text{Tautologies}(\mathcal{L}_m)$, so

$$F \notin \bigcap \{ \text{Tautologies}(\mathcal{L}_n) \mid n \geq 2, n \in \mathbb{N} \}$$

“Fuzzy” logics

$$W = [0, 1]$$

Question: How to define conjunction?

Answer: Desired conditions

$f : [0, 1]^2 \rightarrow [0, 1]$ such that:

- f associative and commutative
- for all $0 \leq A \leq B \leq 1$ and all $0 \leq C \leq 1$ we have $f(A, C) \leq f(B, C)$
- for all $0 \leq C \leq 1$ we have $f(C, 1) = C$.

Definition A function with the properties above is called a t-norm.

Examples of t-norms

Gödel t-norm $f_G(x, y) = \min(x, y)$

Łukasiewicz t-norm $f_L(x, y) = \max(0, x + y - 1)$

Product t-norm $f_P(x, y) = x \cdot y$

Left-continuous t-norm

Definition. A t-norm f is **left-continuous** if for every $x, y \in [0, 1]$ and every sequence $\{x_n\}_{n \in \mathbb{N}}$ with $0 \leq x_n \leq x$ and $\lim_{n \rightarrow \infty} x_n = x$ we have $\lim_{n \rightarrow \infty} f(x_n, y) = f(x, y)$.

Left-continuous t-norm

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The following t-norms are left continuous:

Gödel t-norm $f_G(x, y) = \min(x, y)$

Łukasiewicz t-norm $f_L(x, y) = \max(0, x + y - 1)$

Product t-norm $f_P(x, y) = x \cdot y$

Left continuous t-norms

With every left continuous t-norm f we can associate the following operations:

- $x \circ_f y = f(x, y)$
- $x \oplus_f y = 1 - f(1 - x, 1 - y)$
- $x \Rightarrow_f y = \max\{z \mid f(x, z) \leq y\}$
- $\neg_f x = x \Rightarrow_f 0$

Remark: Left continuity ensures that $\max\{z \mid f(x, z) \leq y\}$ exists.

Validity: $D = \{1\}$

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- $\neg_f x = x \Rightarrow_f 0$

Łukasiewicz t-norm

$$x \circ_{\mathbf{L}} y = \max(0, x + y - 1)$$

$$x \oplus_{\mathbf{L}} y = 1 - \max(0, 1 - x - y)$$

$$x \Rightarrow_f y = \min(1, 1 - x + y)$$

$$\neg x = \min(1, 1 - x) = 1 - x$$

Left continuous t-norms

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- $x \circ_f y = f(x, y)$
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- $\neg_f x = x \Rightarrow_f 0$

Łukasiewicz t-norm

$$x \circ_{\perp} y = \max(0, x + y - 1)$$

$$x \oplus_{\perp} y = 1 - \max(0, 1 - x - y)$$

$$x \Rightarrow_f y = \min(1, 1 - x + y)$$

$$\neg x = \min(1, 1 - x) = 1 - x$$

$$x \wedge_{\perp} y = x \circ_{\perp} (x \Rightarrow y)$$

$$x \vee_{\perp} y = \neg_{\perp}((\neg_{\perp} x) \wedge_{\perp} (\neg_{\perp} y))$$

Left continuous t-norms

With every left continuous t-norm f we can associate the following operations:

- $x \circ_f y = f(x, y)$
- $x \oplus_f y = 1 - f(1 - x, 1 - y)$
- $x \Rightarrow_f y = \max\{z \mid f(x, z) \leq y\}$
- $\neg_f x = x \Rightarrow_f 0$

Gödel t-norm

$$x \circ_G y = \min(x, y)$$

$$x \oplus_G y = \max(x, y)$$

$$x \Rightarrow_G y = \max\{z \mid x \wedge z \leq y\} = \begin{cases} 1 & \text{if } x \leq y \\ y & \text{if } x > y \end{cases}$$

$$\neg_G x = \max\{z \mid x \wedge z = 0\} = \begin{cases} 1 & \text{if } x = 0 \\ 0 & \text{if } x > 0 \end{cases}$$