## Universität Koblenz-Landau

## FB 4 Informatik

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## Exercises for "Non-Classical Logics" <br> Exercise sheet 1

## Exercise 1.1: (2 P)

Determine which of the following formulas are valid/satisfiable/unsatisfiable:
(1) $(P \wedge Q) \rightarrow P$
(2) $Q \wedge \neg Q$
(3) $\neg(\neg P \vee \neg \neg P)$
(4) $((Q \rightarrow P) \wedge(R \rightarrow \neg P)) \rightarrow(\neg Q \vee \neg R)$

## Exercise 1.2: (2 P)

Prove Prop. 1.3 (2): If $N$ is a set of propositional formulas, then $N \models F$ if and only if $N \cup\{\neg F\}$ is unsatisfiable.

Definition: A set of propositional formulas is unsatisfiable, if and only if for every valuation $\mathcal{A}$ there is a formula $G$ in the set such that $\mathcal{A} \not \vDash G$ (i.e. if and only if there is no valuation $\mathcal{A}$ such that $\mathcal{A} \models G$ for all formulae $G$ in the set).

Exercise 1.3: (3 $P$ )
Let $F$ be the following formula:

$$
\neg[((Q \wedge \neg P) \wedge \neg(Q \wedge R)) \rightarrow(Q \wedge \neg P)] \wedge(P \vee R)
$$

(1) Compute the negation normal form (NNF) $F^{\prime}$ of $F$.
(2) Convert $F^{\prime}$ to CNF using:
(a) distributivity of disjunctions over conjunctions?
(b) the satisfiability-preserving transformation described in the lecture.

Exercise 1.4: (2 P)
Consider the formulae $F_{n}=\bigvee_{i=1}^{n}\left(Q_{i} \wedge R_{i}\right)$ for $n \in \mathbb{N}$.
As a function of $n$, how many clauses are in:
(1) the CNF formula $F^{\prime}$ constructed using the distributivity of disjunctions over conjunctions?
(2) the CNF formula $F^{\prime \prime}$ obtained using the satisfiability-preserving translation to clause form?
(3) For which $n$ is the first approach better?

Exercise 1.5: (2 P)
Use the resolution calculus to prove that the following set of clauses is unsatisfiable:

| (1) | $\neg P \vee \neg Q \vee R$ |
| :---: | :---: |
| (2) | $\neg P \vee \neg Q \vee S$ |
| (3) | $P$ |
| (4) | $\neg S \vee \neg R$ |
| (5) | $Q$ |

Exercise 1.6: (2 P)
Assume $S \succ P \succ Q \succ R$. Let $N$ be the following set of clauses:
(1) $\quad \neg Q \vee \neg P$
(2) $\quad R \vee P$
(3) $Q \vee S$
(4) $\quad \neg Q \vee \neg S$

How are the clauses in $N$ ordered w.r.t. the multiset extension of $\succ$ ?

Please submit your solution until Wednesday, October 30, 2013, at 10:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name(s) on your solution.
Submission possibilities:

- By e-mail to sofronie@uni-koblenz.de with the keyword "Homework Non-Classical Logics" in the subject.
- Put it in the box in front of Room B 222.

