

### Exercises for “Non-Classical Logics” Exercise sheet 11

**Exercise 11.1:** (4 P)

Prove that the following formulae are valid using the tableau calculus presented in the lecture.

- (1)  $\diamond(P \vee Q) \rightarrow (\diamond P \vee \diamond Q)$
- (2)  $(\diamond P \vee \diamond Q) \rightarrow \diamond(P \vee Q)$
- (3)  $\diamond(P \wedge Q) \rightarrow (\diamond P \wedge \diamond Q)$

**Exercise 11.2:** (2 P)

Prove that the formula  $A$  is satisfiable using the tableau calculus presented in the lecture:

$$A : \neg((\diamond P \wedge \diamond Q) \rightarrow \diamond(P \wedge Q))$$

and construct a Kripke model  $\mathcal{K} = (S, R, I)$  and a state  $s \in S$  such that  $(\mathcal{K}, s) \models A$  using a saturated tableau for  $A$ .

**Exercise 11.3:** (2 P)

Construct a saturated or closed tableau starting from the following prefixed formula:

$$T((\Box \diamond P \wedge \diamond P) \rightarrow \diamond \Box P)$$

**Exercise 11.4:** (2 P)

Compute the translation into first order logic used for checking the validity of a modal formula  $\Phi$  (of the form  $\exists x P_{\neg\Phi}(x) \wedge \text{Rename}(\neg\Phi)$ ) for the following formulae:

- (1)  $\Phi_1 : \diamond(P \wedge Q) \rightarrow (\diamond P \wedge \diamond Q)$
- (2)  $\Phi_2 : (\diamond P \vee \diamond Q) \rightarrow \diamond(P \vee Q)$
- (3)  $\Phi_3 : ((\Box \diamond P \wedge \diamond P) \rightarrow \diamond \Box P)$

Please submit your solution until Tuesday, January 21, 2014 at 16:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name(s) on your solution.

Submission possibilities:

- By e-mail to [sofronie@uni-koblenz.de](mailto:sofronie@uni-koblenz.de) with the keyword “Homework Non-Classical Logics” in the subject.
- Put it in the box in front of Room B 222.