# Universität Koblenz-Landau 

## FB 4 Informatik

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December 3, 2014

## Exercises for "Non-Classical Logics" <br> Exercise sheet 6

Exercise 6.1: (4 P)
Let $g:\{0, u, 1\} \times\{0, u, 1\} \rightarrow\{0, u, 1\}$ be the function defined as described in the table:

| $g$ | 0 | $u$ | 1 |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | $u$ |
| $u$ | 0 | $u$ | 0 |
| 1 | $u$ | 0 | 1 |

Construct a formula $F$ with propositional variables $P_{1}, P_{2}$ in the (functionally complete) propositional logic $\mathcal{L}_{3}^{+}$defined in the lecture from 3.12 .2014 which "has the same truth table as $g "$, i.e. has the property that for every $\mathcal{A}:\left\{P_{1}, P_{2}\right\} \rightarrow\{0, u, 1\}:$

$$
g\left(\mathcal{A}\left(P_{1}\right), \mathcal{A}\left(P_{2}\right)\right)=\mathcal{A}(F)
$$

Exercise 6.2: ( 6 P)
Use the semantic tableau calculus for the many-valued logic $\mathcal{L}_{3}$ to prove that the following formulae are $\mathcal{L}_{3}$ tautologies:
(1) $\neg \neg A$ id $A$
(2) $\neg(A \vee B)$ id $(\neg A \wedge \neg B)$
(3) $\sim(\exists x F(x))$ id $\forall x(\sim F(x))$
(Hint: To avoid the problem of having to use the definition of the operator id the truth table of id can be used for devising suitable expansion rules for the tableau calculus.)

[^0]- By e-mail to sofronie@uni-koblenz.de with the keyword "Homework Non-Classical Logics" in the subject.
- Put it in the box in front of Room B 222.


[^0]:    Please submit your solution until Monday, December 8, 2014, at 18:00. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name(s) on your solution.

    Submission possibilities:

