## Universität Koblenz-Landau

## FB 4 Informatik

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## Exercises for <br> "Advances in Theoretical Computer Science" <br> Exercise sheet 3

## Exercise 3.1:

Let $L_{1}, L_{2}, L_{3}$ be languages, where $L_{2}$ is recursively enumerable and $L_{3}$ is decidable. Prove or refute the following statements:

1. If $L_{1} \subseteq L_{3}$, then $L_{1}$ is decidable.

2 . If $L_{3} \subseteq L_{1}$, then $L_{1}$ is decidable.
3 . If $L_{1} \subseteq L_{2}$, then $L_{1}$ is recursively enumerable.
4. If $L_{2} \subseteq L_{1}$, then $L_{1}$ is recursively enumerable.

## Exercise 3.2:

Decide for each of the following problems if it is decidable or not. Justify your answer.
a) $P_{1}:=\left\{n \in \mathbb{N} \mid \mathcal{M}_{n}\right.$ does not hold on empty input $\}$
b) $P_{2}:=\left\{n \in \mathbb{N} \mid L\left(\mathcal{M}_{n}\right)=\emptyset\right\}$
c) $P_{3}:=\left\{(m, n) \in \mathbb{N} \times \mathbb{N} \mid L\left(\mathcal{M}_{m}\right) \cap L\left(\mathcal{M}_{n}\right)=\emptyset\right\}$
d) $P_{4}:=\left\{(m, n) \in \mathbb{N} \times \mathbb{N} \mid L\left(\mathcal{M}_{m}\right) \subseteq L\left(\mathcal{M}_{n}\right)\right\}$
e) $P_{5}:=\left\{(n, w) \in \mathbb{N} \times \Sigma^{*} \mid\right.$ For input $w, \mathcal{M}_{n}$ does not reach another configuration after $s, \# w \#$ where the head is on a blank $(\#)\}$
f) $P_{6}:=\left\{(n, w, s) \in \mathbb{N} \times \Sigma^{*} \times \mathbb{N} \mid \mathcal{M}_{n}\right.$ halts on input $w$ after at most $s$ steps $\}$

Remarks:

- $\mathcal{M}_{n}$ denotes the Turing machine with Gödel number $n$.
- $L(\mathcal{M})$ is the language accepted by the Turing machine $\mathcal{M}$ (i.e. the set of all words accepted by $\mathcal{M})$.

Hint: To prove undecidability you can for instance use properties of decidable languages (e.g. the fact that the complement of a decidable language is decidable) or a reduction to a problem which was already proved to be undecidable: (1) You are allowed to use the undecidability of the halting problem $H A L T$, of $H_{0}$ or of $K$ (notation as in the lecture); (2) if you have proven the (un-)decidability of $P_{i}$, you may use this result for any of the next tasks.

The submission of the solutions is not compulsory. If you want to submit your solutions, please do so until Tuesday, 12.11.2013, 10:00 s.t.. Joint solutions prepared by up to three persons are allowed. Please do not forget to write your name on your solution.
Submission possibilities:

- By e-mail to mbender@uni-koblenz.de with the keyword "Homework ACTCS" in the subject.
- Put it in the box in front of Room B 222.

