## **Advanced Topics in Theoretical Computer Science**

Part 2: Register machines

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- Recapitulation: Turing machines and Turing computability
- Register machines (LOOP, WHILE, GOTO)
- Recursive functions
- The Church-Turing Thesis
- Computability and (Un-)decidability
- Complexity

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- LOOP Programs
- WHILE Programs
- GOTO Programs
- Relationships between LOOP, WHILE, GOTO
- Relationships between register machines and Turing machines

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The register machine gets its name from its one or more "registers":

In place of a Turing machine's tape and head (or tapes and heads) the model uses multiple, uniquely-addressed registers, each of which holds a single positive integer.

#### In comparison to Turing machines:

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similar to ...
the imperative kernel of programming languages
pseudo-code

Computation of a mod b (pseudocode)

```
r := a;
while r \ge b do
r := r - b
end;
return r
```

#### **Definition: Questions**

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- Data structures
  - Unbounded but finite number of registers denoted  $x_1, x_2, x_3, \ldots, x_n$ ; each register contains a natural number (no arrays, objects, ...)

- Atomic instructions:
  - Increment/Decrement a register

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  - Increment/Decrement a register
- Input/Output
  - Input: n input values in the first n registers
     All the other registers are 0 at the beginning.
  - **Output:** In register n + 1.

#### **Syntax**

#### **Definition**

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## **Example: WHILE Programs**

#### **Syntax**

#### **Definition**

- Atomic programs: For each register  $x_i$ :
  - $x_i := x_i + 1$
  - $x_i := x_i 1$

are WHILE instructions and also WHILE programs.

- If  $P_1$ ,  $P_2$  are WHILE programs then
  - $P_1$ ;  $P_2$  is a WHILE program
- If *P* is a WHILE program then
  - while  $x_i \neq 0$  do P end is a WHILE program (and a WHILE instruction)

Syntax Indexes (numbers for the lines in the program)  $j \ge 0$ 

#### **Definition**

- Atomic programs:
  - $x_i := x_i + 1$
  - $x_i := x_i 1$

are GOTO instructions for each register  $x_i$ .

- If  $x_i$  is a register and j is an index then
  - if  $x_i = 0$  goto j is a GOTO instruction.
- If  $I_1, \ldots, I_k$  are GOTO instructions and  $j_1, \ldots, j_k$  are indices then
  - $-j_1:I_1;\ldots;j_k:I_k$  is a GOTO program

#### **Definition**

A register machine is a machine consisting of the following elements:

- A finite (but unbounded) number of registers  $x_1, x_2, x_3, \dots, x_n$ ; each register contains a natural number.
- A LOOP-, WHILE- or GOTO-program.

# Register Machines: State

#### **Definition** (State of a register machine)

The state s of a register machine is a map:

$$s: \{x_i \mid i \in \mathbb{N}\} \to \mathbb{N}$$

which associates with every register a natural number as value.

# Register Machines: State

#### **Definition (Initial state; Input)**

Let  $m_1, \ldots, m_k \in \mathbb{N}$  be given as input to a register machine.

In the input state  $s_0$  we have

- $s_0(x_i) = m_i$  for all  $1 \le i \le k$
- $s_0(x_i) = 0$  for all i > k

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#### **Definition (Output)**

If a register machine started with the input  $m_1, \ldots, m_k \in \mathbb{N}$  halts in a state  $s_{\text{term}}$  then:

$$s_{\text{term}}(x_{k+1})$$

is the output of the machine.